



# **Token-Based Adaptive MAC for Two-Hop Device-to-Device Communications**

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# Device-to-Device (D2D) Communications

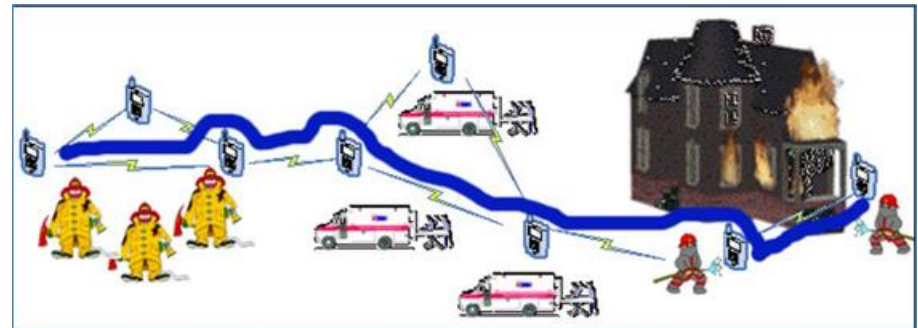
## ■ Definition

- Self-organized nodes interconnected for communications in a distributed manner
  - Low cost and simplified implementation
- Ad hoc networking
  - Smart wireless devices (strong computing and smart sensing capabilities)



## ■ Characteristics

- Infrastructure-less
- Node movement
- Multi-hop communications
- Increasing node number



## ■ Applications

- Emergency communications for disaster recovery
- Smarting sensing for flood monitoring

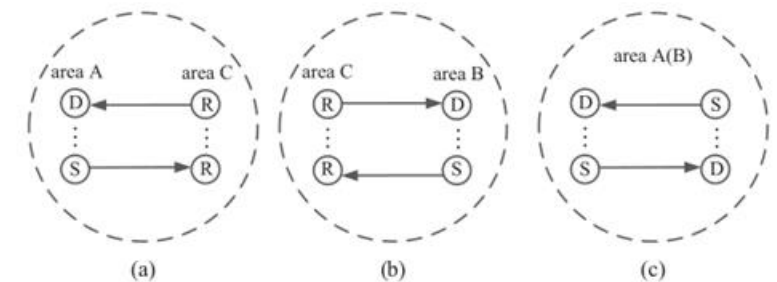
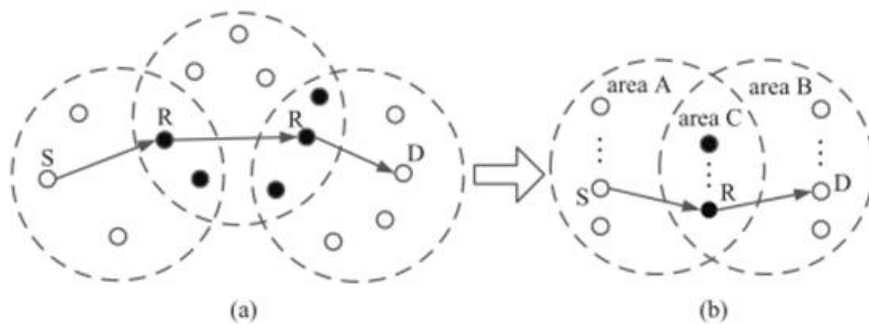
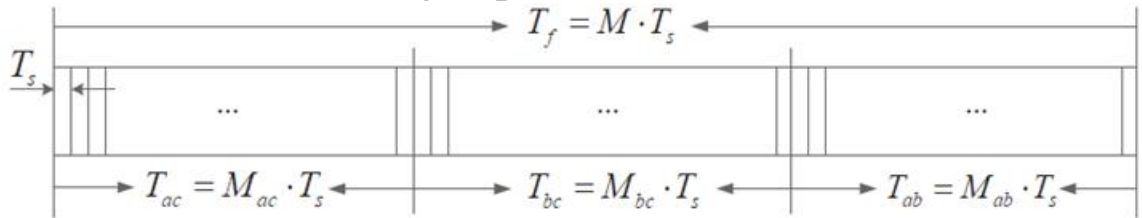
# Medium Access Control

- **Medium access control (MAC)**
  - Coordinate nodes' channel access
- **Objective**
  - Low end-to-end packet transmission delay
- **Challenges for MAC in D2D communications**
  - Distributed network operation
    - Distributed MAC
  - Varying number of users
    - Traffic-adaptive MAC
  - Enlarged network region with an increasing number of users
    - Interference-aware and scalable MAC



# System Model

- A two-hop network with an error free channel
- $N_a, N_b, N_c$  in areas A, B and C
- Sensing data traffic
- MAC superframe structure
  - Spatial reservation of time slots for different node groups
- Four one-hop subnetworks



(a) A general multi-hop MANET. (b) A simplified two-hop network.

Packet transmissions for four one-hop subnetworks during (a)  $T_{ac}$ , (b)  $T_{bc}$ , (c)  $T_{ab}$ .

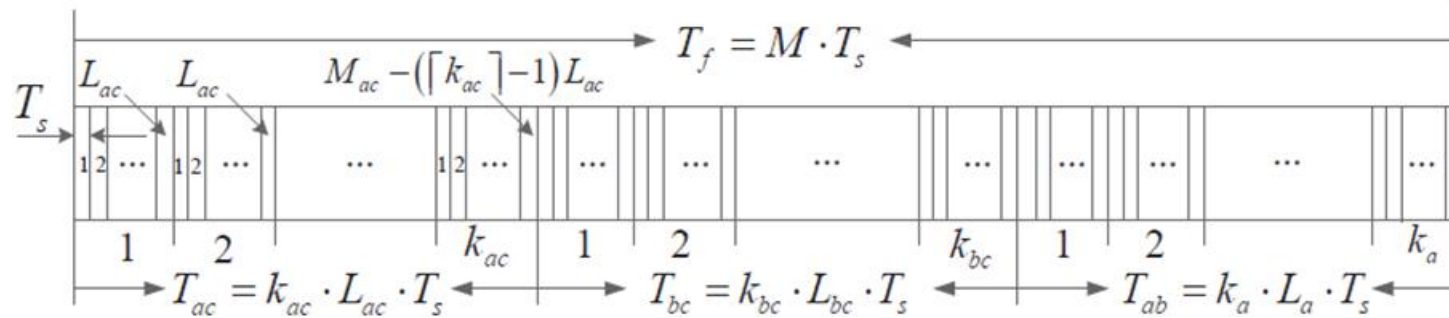
# Token-based Adaptive MAC

- **Probabilistic token passing scheme**

- Four token rings  $R_{ac}$ ,  $R_{bc}$  and  $R_a$  ( $R_b$ ) formed among each subnetwork
- Tokens circulated probabilistically within each node group
  - One complete token rotation cycle
- Numbers of token rotation cycles  $k_{ac}$ ,  $k_{bc}$  and  $k_a$  ( $k_b$ ) in durations  $T_{ac}$ ,  $T_{bc}$  and  $T_{ab}$ .

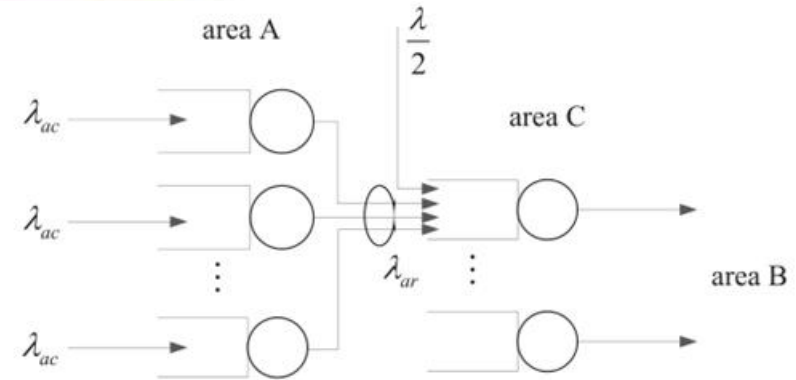
- **Important MAC parameters**

- Number of token rotation cycles  $k_j$  for each node group
- Total number of time slots  $M$



# Performance Analysis

- **Compound packet arrival rate**
  - Poisson traffic approximation on each relay node (an extension to Kleinrock independence approximation)



- **Average end-to-end delay**

$$D_{ab} = \sum_{(n,j) \in \{(ac,ac), (cb,bc)\}} \left( \frac{\varepsilon_j}{k_j} + \frac{\lambda_n [\alpha_j k_j^2 + \beta_j k_j + \gamma_j]}{2(k_j - \lambda_n \varepsilon_j)} \right) \quad D_{ba} = \sum_{(n,j) \in \{(bc,bc), (ca,ac)\}} \left( \frac{\varepsilon_j}{k_j} + \frac{\lambda_n [\alpha_j k_j^2 + \beta_j k_j + \gamma_j]}{2(k_j - \lambda_n \varepsilon_j)} \right)$$

where  $\alpha_j = L_j^2$ ;  $\beta_j = -\frac{5L_j^2 + 12ML_j + 1}{6}$ ;  $\gamma_j = M^2 + 2ML_j$ ;  $\varepsilon_j = M$

- **Average delay for local transmissions in areas A and B**

$$D_j = \frac{\varepsilon_j}{k_j} + \frac{\lambda_j [\alpha_j k_j^2 + \beta_j k_j + \gamma_j]}{2(k_j - \lambda_j \varepsilon_j)} \quad (j = a, b)$$

where  $\alpha_j = L_j^2$ ;  $\beta_j = -\frac{5L_j^2 + 12ML_j + 1}{6}$ ;  $\gamma_j = M^2 + 2ML_j$ ;  $\varepsilon_j = M$

# Optimal MAC Parameters

- Average end-to-end delay minimization with a given  $M$**

$$(P1) : \min_{\mathbf{k}=[k_{ac}, k_{bc}, k_a, k_b]} \{\max\{D_{ab}(k_{ac}, k_{bc}), D_{ba}(k_{ac}, k_{bc})\}\}$$

$$\text{s.t.} \begin{cases} k_{ac}L_{ac} + k_{bc}L_{bc} + k_aL_a = M \\ k_aL_a = k_bL_b \\ \rho_n = \frac{\lambda_n}{\mu_j} < 1 \quad (n, j) \in \{(ca, ac), (cb, bc)\} \\ \rho_j = \frac{\lambda_j}{\mu_j} < 1 \quad (j = a, b) \\ D_j(k_j) \leq D_{th} \quad (j = a, b) \\ k_j \geq 1 \quad (j = ac, bc, a, b). \end{cases}$$



**Solution: Problem decoupling**



$$(SP1) : \min_{\mathbf{k}_1=[k_a, k_b]} k_aL_a \quad (SP2) : \min_{\mathbf{k}_2=[k_{ac}, k_{bc}]} \{\max\{D_{ab}(k_{ac}, k_{bc}), D_{ba}(k_{ac}, k_{bc})\}\}$$

$$\text{s.t.} \begin{cases} k_aL_a = k_bL_b \\ \rho_j < 1 \quad (j = a, b) \\ D_j(k_j) \leq D_{th} \quad (j = a, b) \\ k_j \geq 1 \quad (j = a, b). \end{cases}$$

$$\text{s.t.} \begin{cases} k_{ac}L_{ac} + k_{bc}L_{bc} = M^* \\ \rho_n < 1 \quad (n = ca, cb) \\ k_j \geq 1 \quad (j = ac, bc) \end{cases}$$

$$\text{where } M^* = M - k_a^*L_a$$

A convex subproblem

A biconvex subproblem



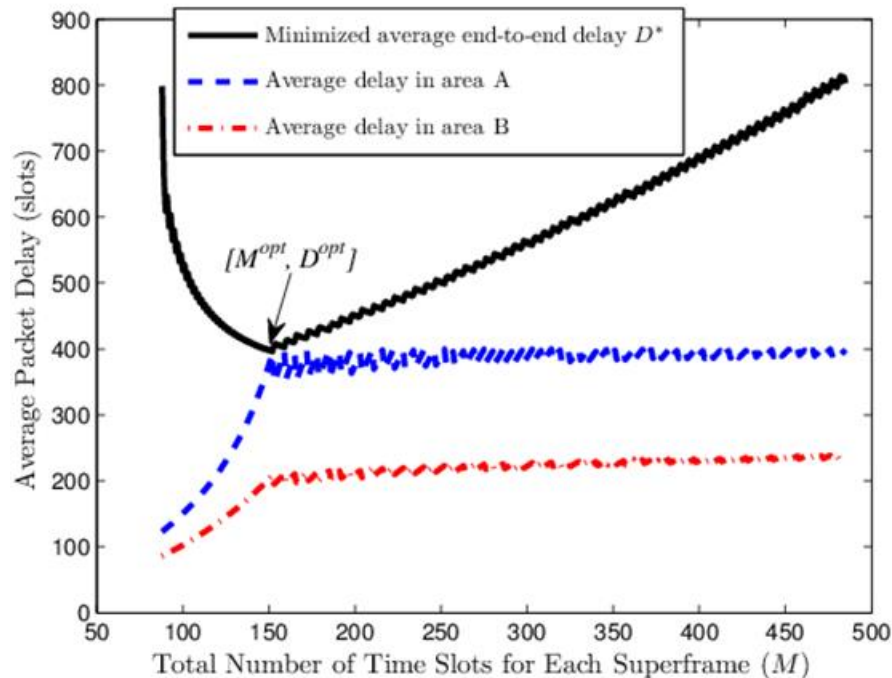
$$\mathbf{k}^* = [k_{ac}^*, k_{bc}^*, k_a^*, k_b^*].$$

- Calculate optimal total number of time slots**
  - Optimal parameter set  $[k_j^{opt}, M^{opt}]$

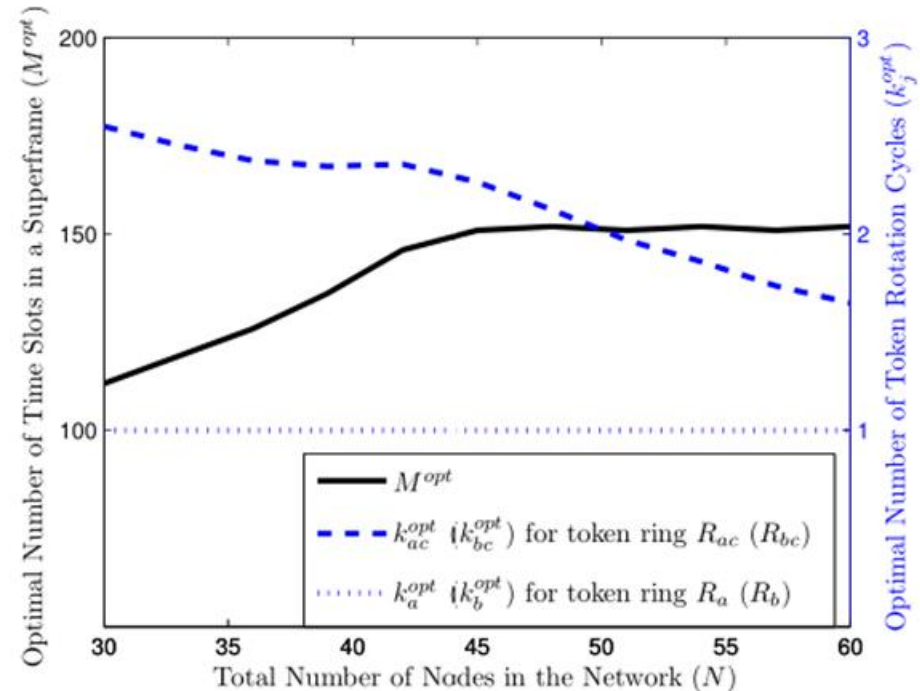
# Numerical Results (1)

## Optimal MAC parameters

- Optimal total number of time slots  $M^{opt}$  (slot duration: 1ms)
- Optimal number of token rotation cycles  $k_j^{opt}$



$$N_a = 20, N_b = 15, N_c = 15$$



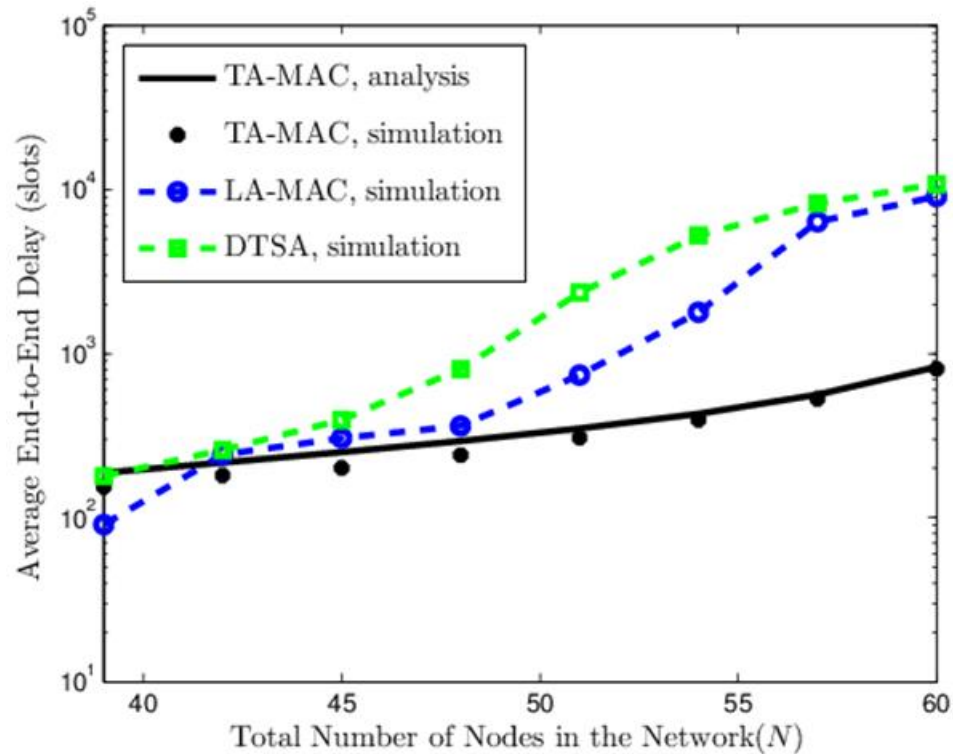
$$N_a = N_b = N_c$$



# Numerical Results (2)

## ■ Performance comparison

- Load adaptive MAC (LA-MAC) and Dynamic time slot allocation (DTSA)



# Conclusions

- **A scalable token-based adaptive MAC for a two-hop network**
  - A TDMA-based superframe structure for spatial reservation of time resources among different groups of nodes
  - Probabilistic token passing for distributed time slot allocation
  - Closed-form performance analytical modeling for average end-to-end packet transmission delay in terms of MAC parameters (i.e., numbers of token rotation cycles and the superframe length)
  - Determine optimal MAC parameters to achieve minimal average end-to-end delay

# Future Plan

- Specifically, we will investigate
  - how to form **multiple concurrent device-to-device groups** that share devices of richer power as cluster heads;
  - how to incorporate **various types of data transmission** to CAFFEWS data server and enhance transmission reliability;
  - how to improve initial source selection for **data dissemination via device-to-device communications**;
  - how to take advantage of user-provided information to expand the scope of data dissemination; and
  - to collaborate with other researchers in Theme 3 to refine/validate/test CAFFEWS.

Thank you!